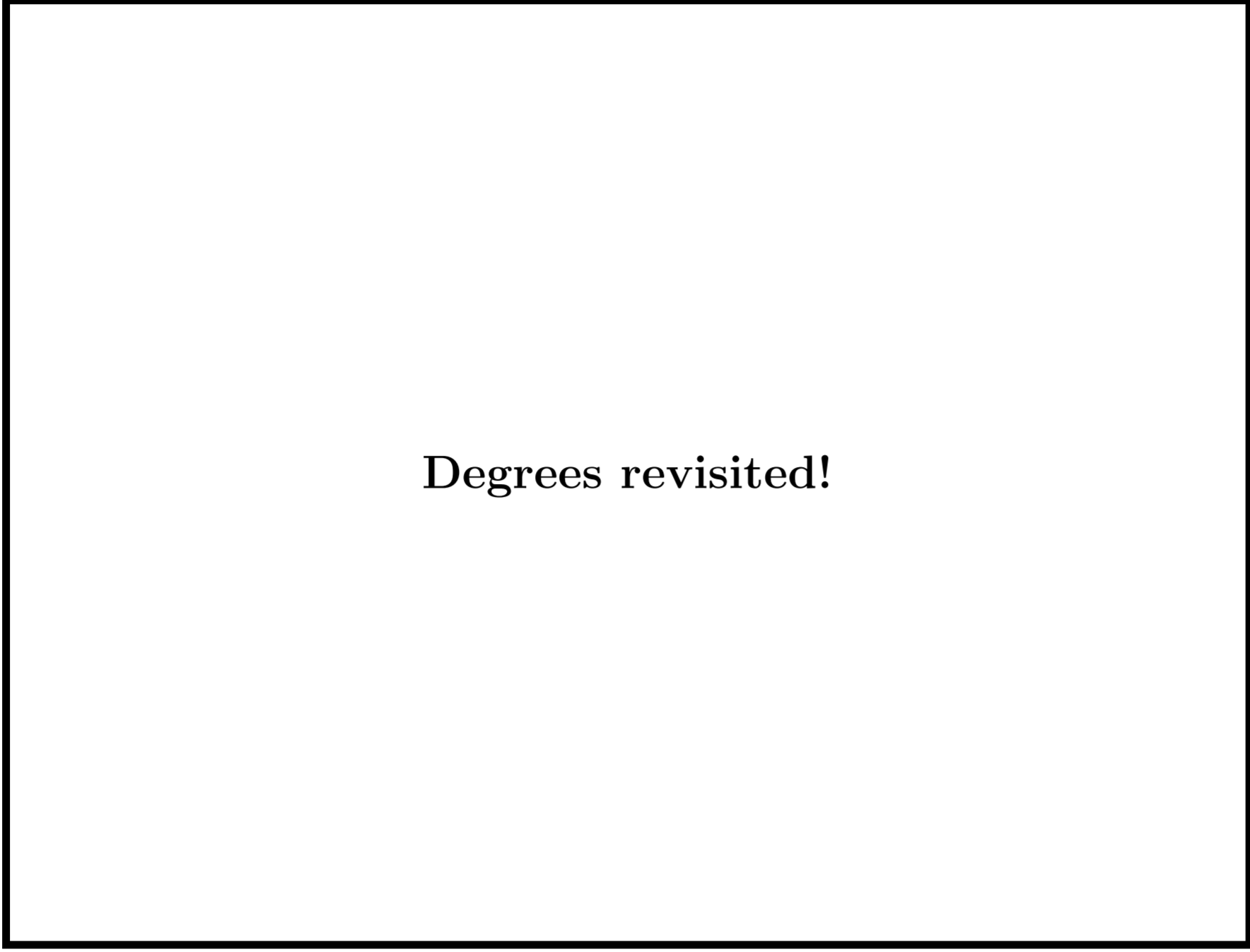


**Adventures in random graphs:
Models, structures and algorithms**

Armand M. Makowski

**ECE & ISR/HyNet
University of Maryland at College Park
armand@isr.umd.edu**

LECTURE 4
Scale-free networks



Degrees revisited!

Consider a directed graph $G = (V, E)$ with $V = \{1, \dots, n\}$ with link assignments $\{\chi_{ij}, 1 \leq i, j \leq n\}$.

The **out-degree** of node i :

$$d_i^+ := \sum_{j=1, j \neq i}^n \chi_{ij}$$

The **in-degree** of node i :

$$d_i^- := \sum_{j=1, j \neq i}^n \chi_{ji}$$

Basic identity

$$\sum_{i=1}^n d_i^+ = \sum_{i=1}^n d_i^- = |E|$$

For undirected graphs

For **undirected** graphs, the **degree** of node i is simply

$$d_i = d_i^+ = d_i^-, \quad i \in V$$

Basic identity

$$\sum_{i=1}^n d_i = |E|$$

or

$$\sum_{i=1}^n d_i = 2 \cdot \text{Total number of **undirected** edges in } G$$

Empirical degree distribution in $G = (V, E)$

$$\begin{aligned} P(d) &= \frac{1}{n} \sum_{i=1}^n \mathbf{1}[d_i = d] \\ &= \frac{|\{i = 1, \dots, n : d_i = d\}|}{n}, \quad d = 0, 1, \dots, n-1 \end{aligned}$$

so that

$$d_{\text{Avg}} = \sum_{d=0}^{n-1} d \cdot P(d) = \frac{1}{n} \sum_{i=1}^n d_i$$

No random graph model assumed – Data driven

The random node viewpoint

With ν denoting a rv **uniformly** distributed over the index set $\{1, \dots, n\}$, define the **rv** given by

$$D = d_\nu$$

It is plain that

$$\begin{aligned} \mathbb{P}[D = d] &= \sum_{i=1}^n \mathbb{P}[d_i = d, \nu = i] \\ &= \sum_{i=1}^n \mathbf{1}[d_i = d] \mathbb{P}[\nu = i] = P(d) \end{aligned}$$

Random sampling interpretation of $P(\cdot)$!

The random edge viewpoint

Let ε be a rv **uniformly** distributed over the edge set E of V , and write $\varepsilon = (i, j)$ with $1 \leq i \neq j \leq n$. Next, **independently** of the outcome, select either node with probability $1/2$, and denote the resulting node by μ . Set

$$D^* = d_\mu$$

Can we compute the pmf

$$P^*(d) = \mathbb{P}[D^* = d], \quad d = 0, \dots, n-1?$$

For each $d = 0, 1, \dots, n - 1$,

$$\begin{aligned}\mathbb{P}[D^* = d] &= \mathbb{P}[d_\mu = d] \\ &= \sum_{k=1}^n \mathbb{P}[\mu = k, d_\mu = d] \\ &= \sum_{k=1}^n \mathbb{P}[\mu = k, d_k = d] \\ &= \sum_{k=1}^n \mathbf{1}[d_k = d] \mathbb{P}[\mu = k]\end{aligned}$$

By construction,

$$\mathbb{P}[\mu = k] = \frac{\frac{1}{2}d_+(k) + \frac{1}{2}d_-(k)}{|E|} = \frac{d(k)}{|E|}$$

so that

$$\begin{aligned} \mathbb{P}[D^* = d] &= \sum_{k=1}^n \mathbf{1}[d_k = d] \frac{d(k)}{|E|} \\ &= \frac{d}{|E|} \cdot \sum_{k=1}^n \mathbf{1}[d_k = d] \\ &= \frac{d}{|E|} \cdot nP(d) \end{aligned}$$

$$P^*(d) = \frac{n}{|E|} \cdot d P(d), \quad d = 0, 1, \dots, n-1$$

Note that

$$\sum_{d=0}^{n-1} P^*(d) = 1$$

since

$$\begin{aligned} \sum_{d=0}^{n-1} d P(d) &= \sum_{d=0}^{n-1} d \cdot \left(\frac{1}{n} \sum_{i=1}^n \mathbf{1}[d_i = d] \right) \\ &= \frac{1}{n} \sum_{i=1}^n \sum_{d=0}^{n-1} d \cdot \mathbf{1}[d_i = d] \\ &= \frac{1}{n} \sum_{i=1}^n d_i \\ &= \frac{|E|}{n} \end{aligned}$$

Arbitrary degree distributions?

Degrees in undirected graphs

Is it always the case that non-negative integers $(d_i, i \in V)$ are the degree sequence of a graph $G = (V, E)$?

Erdős and Gallai (1960): The positive integers d_1, \dots, d_n are the degrees of a simple graph with n vertices if and only if

$$d_1 + \dots + d_n \text{ even}$$

and

$$\sum_{j=1}^k d_{(j)} \leq k(k-1) + \sum_{\ell=k+1}^n \min(k, d_{(\ell)}), \quad k = 1, \dots, n$$

with $d_{(1)}, \dots, d_{(n)}$ denoting d_1, \dots, d_n in decreasing order.

A simple fact

Theorem 1 *With i.i.d.. \mathbb{N}_0 -valued rvs $\{D, D_n, n = 1, 2, \dots\}$, we have*

$$\lim_{n \rightarrow \infty} \mathbb{P}[D_1 + \dots + D_n \text{ even}] = \frac{1}{2}$$

provided

$$0 < \mathbb{P}[D \text{ even}] < 1$$

Proof: For each $n = 1, 2, \dots$, note that

$$\mathbb{E}[(-1)^{D_1 + \dots + D_n}] = \mathbb{E}[(-1)^D]^n = (2\mathbb{P}[D \text{ even}] - 1)^n$$

while

$$\mathbb{E}[(-1)^{D_1 + \dots + D_n}] = 2\mathbb{P}[D_1 + \dots + D_n \text{ even}] - 1$$

In many deployed networks, data analysis shows a Pareto-like behavior, e.g.,

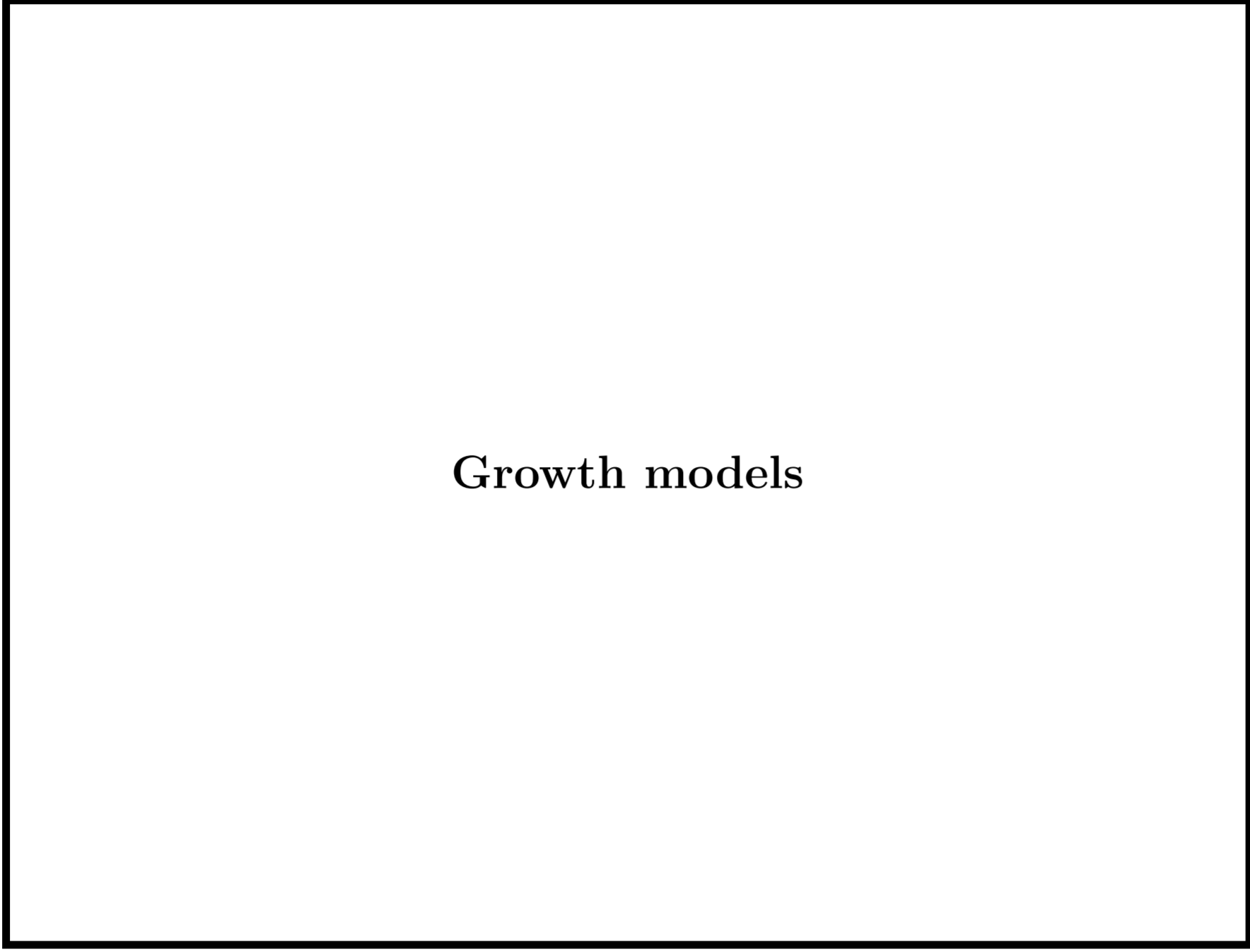
$$P(d) \sim d^{-\gamma}, \quad d = 1, \dots, n-1$$

for some $\gamma > 2$.

Why generating models that capture this property?

- From local rules to global properties!
 - Impact on the behavior of dynamic processes on such networks
-

- Growth models
- Optimization-based models
- Threshold random graphs
- !!!



Growth models

What are growth models?

Sequence of **random** graph models

$$\{\mathbb{G}_t, t = 0, 1, \dots\}$$

with rules/mechanisms

$$V_{t+1} \leftarrow V_t$$

and

$$\mathbb{G}_{t+1} \leftarrow (\mathbb{G}_t, V_{t+1})$$

Usually **Markov**-like random mechanisms with the requirements that

$$V_t \subseteq V_{t+1}, \quad t = 0, 1, \dots$$

and

$$\mathbb{G}_t = \text{Subgraph of } \mathbb{G}_{t+1}, \quad t = 0, 1, \dots$$

Two basic rules

- Preferential attachment – “The rich get richer”
- Copying

Very long history with many more models recently proposed:

- Yule (1925)
- Barabási and Albert (1999)
- Kleinberg et al.
-

Key features

- Nodes are non-homogeneous – Infinite node population
- Yet not a dynamical network model
- No node deletion and no edge deletion!
- Interest in asymptotic behavior of the sample degree distributions

$$\frac{1}{|V_t|} \sum_{i \in V_t} \mathbf{1}[D_t(i) = d], \quad \begin{array}{l} d = 0, 1, \dots \\ t = 0, 1, \dots \end{array}$$

where

$$D_t(i) = \text{Degree of node } i \text{ in } \mathbb{G}_t$$

The BA model (Simplified version)

We are given

- A connected graph $G = (V, E)$ with

$$V = \{-(K-1), \dots, -1, 0\}$$

for some integer K in \mathbb{N}

- An **attachment** function, i.e., a monotone increasing mapping $\varphi : \mathbb{N}_0 \rightarrow \mathbb{R}_+$ with $\varphi(u) = 0$ only if $u = 0$, i.e.,

$$\varphi(u) = u^\nu \quad (\nu > 0) \quad \text{or} \quad \varphi(u) = e^{\lambda u} \quad (\lambda > 0)$$

Recursively

For $t = 0$:

$$\mathbb{G}_0 = G$$

For $t = 0, 1, \dots$: With $\mathbb{G}_t = (V_t, \mathbb{E}_t)$ available, we construct $\mathbb{G}_{t+1} = (V_{t+1}, \mathbb{E}_{t+1})$ as follows:

- A new **node**, labelled $t + 1$, arrives so that

$$V_{t+1} = V_t \cup \{t + 1\}$$

- A new **edge** is created, linking node $t + 1$ to one of the existing nodes, say X_{t+1} , in V_t so that

$$\mathbb{E}_{t+1} = \mathbb{E}_t \cup \{(X_{t+1}, t + 1)\}$$

The preferential attachment rule

The nodes $\{X_{t+1}, t = 0, 1, \dots\}$ are randomly generated as follows:

At time $t + 1$,

$$\mathbb{P}[X_{t+1} = x | X_1, \dots, X_t] = \frac{\varphi(D_t(x))}{\sum_{y \in V_t} \varphi(D_t(y))}, \quad x \in V_t$$

with

$$D_t(i) = \text{Degree of node } i \text{ in } \mathbb{G}_t, \quad i \in V_t$$

Note that

$$D_t(i) > 0, \quad \begin{array}{l} i = -(K-1), \dots, t \\ t = 0, 1, \dots \end{array}$$

A basic recursion

For each $t = 0, 1, \dots$,

$$D_{t+1}(i) = \begin{cases} D_t(i) + \mathbf{1}[X_{t+1} = i] & \text{if } i \in V_t \\ 1 & \text{if } i = t + 1 \end{cases}$$

with

$$D_0(i) = \text{Degree of node } i \text{ in } G, \quad i \in V_0$$

Obviously, by construction, we have

$$|\mathbb{E}_{t+1}| = |\mathbb{E}_t| + 2, \quad t = 0, 1, \dots$$

so that

$$|\mathbb{E}_t| = |E_0| + 2t, \quad t = 0, 1, \dots$$

Hence

$$\sum_{i \in V_t} D_t(i) = |E_0| + 2t, \quad t = 0, 1, \dots$$

The linear case ($\varphi(u) = u, u \geq 0$)

The nodes $\{X_{t+1}, t = 0, 1, \dots\}$ are now randomly generated as follows: At time $t + 1$,

$$\begin{aligned} \mathbb{P}[X_{t+1} = x | X_1, \dots, X_t] &= \frac{D_t(x)}{\sum_{y \in V_t} D_t(y)} \\ &= \frac{D_t(x)}{|E_0| + 2t}, \quad x \in V_t \end{aligned}$$

For $t = 0, 1, \dots$, define the counting rvs

$$N_t(d) = \sum_{x \in V_{t+1}} \mathbf{1}[D_t(x) = d], \quad d = 1, 2, \dots$$

Note that $N_t(0) = 0$.

For $d = 1$:

$$N_{t+1}(1) = \begin{cases} N_t(1) & \text{if } D_t(X_{t+1}) = 1 \\ N_t(1) + 1 & \text{if } D_t(X_{t+1}) > 1 \end{cases}$$

so that

$$N_{t+1}(1) = N_t(1) + \mathbf{1} [D_t(X_{t+1}) \neq 1]$$

$$\begin{aligned}
\mathbb{E}[N_{t+1}(1)|\mathcal{F}_t] &= N_t(1) + \mathbb{P}[D_t(X_{t+1}) \neq 1|\mathcal{F}_t] \\
&= N_t(1) + 1 - \mathbb{P}[D_t(X_{t+1}) = 1|\mathcal{F}_t] \\
&= N_t(1) + 1 - \sum_{x \in V_t} \mathbf{1}[D_t(x) = 1] \frac{D_t(x)}{|E_0| + 2t} \\
&= N_t(1) + 1 - \sum_{x \in V_t} \mathbf{1}[D_t(x) = 1] \frac{1}{|E_0| + 2t} \\
&= N_t(1) + 1 - \frac{N_t(1)}{|E_0| + 2t} \\
&= N_t(1) \cdot \left(1 - \frac{1}{|E_0| + 2t}\right) + 1
\end{aligned}$$

$$\mathbb{E}[N_{t+1}(1)] = \mathbb{E}[N_t(1)] \cdot \left(1 - \frac{1}{|E_0| + 2t}\right) + 1$$

For $d > 1$:

$$N_{t+1}(d) = N_t(d) + \mathbf{1}[D_t(X_{t+1}) = d - 1] - \mathbf{1}[D_t(X_{t+1}) = d]$$

$$\begin{aligned} \mathbb{E}[N_{t+1}(d)|\mathcal{F}_t] &= N_t(d) + \mathbb{P}[D_t(X_{t+1}) = d - 1|\mathcal{F}_t] - \mathbb{P}[D_t(X_{t+1}) = d|\mathcal{F}_t] \\ &= N_t(d) + \sum_{x \in V_t} \mathbf{1}[D_t(x) = d - 1] \mathbb{P}[X_{t+1} = x|\mathcal{F}_t] \\ &\quad - \sum_{x \in V_t} \mathbf{1}[D_t(x) = d] \mathbb{P}[X_{t+1} = x|\mathcal{F}_t] \\ &= N_t(d) + \sum_{x \in V_t} \mathbf{1}[D_t(x) = d - 1] \frac{d - 1}{|E_0| + 2t} \end{aligned}$$

$$\begin{aligned}
& - \sum_{x \in V_t} \mathbf{1}[D_t(x) = d] \frac{d}{|E_0| + 2t} \\
&= N_t(d) + \frac{d-1}{|E_0| + 2t} \cdot N_t(d-1) - \frac{d}{|E_0| + 2t} \cdot N_t(d) \\
&= N_t(d) \left(1 - \frac{d}{|E_0| + 2t}\right) + \frac{d-1}{|E_0| + 2t} \cdot N_t(d-1)
\end{aligned}$$

$$\begin{aligned}
& \mathbb{E}[N_{t+1}(d)] \\
&= \mathbb{E}[N_t(d)] \left(1 - \frac{d}{|E_0| + 2t}\right) + \frac{d-1}{|E_0| + 2t} \cdot \mathbb{E}[N_t(d-1)]
\end{aligned}$$

Theorem 2 *The limits*

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[N_t(d)]}{|V_t|} = \pi(d), \quad d = 1, 2, \dots$$

all exist and satisfy

$$\pi(d) = \frac{d-1}{d+2} \pi(d-1), \quad d = 2, 3, \dots$$

with $\pi(1) = \frac{2}{3}$.

Easy to see that

$$\pi(d) \sim \frac{4}{d^3}$$

A stronger results is available

Theorem 3 *For each $d = 1, 2, \dots$, we have*

$$\lim_{t \rightarrow \infty} \frac{N_t(d)}{|V_t|} = \pi(d) \quad a.s.$$

with $\pi(d)$ as given earlier.

More general models (I)

With α in $[0, 1)$, the **new** node labeled $t + 1$ is attached to one of the existing nodes, say X_{t+1} in V_t ,

- With probability α , **uniformly** amongst the nodes in V_t
 - With probability $1 - \alpha$, according to a **preferential** attachment mechanism
-

With a linear attachment function, Theorems 2 and 3 still hold with

$$\pi(d) \sim d^{-\gamma} \quad \text{with} \quad \gamma = \frac{3 - \alpha}{1 - \alpha}$$

More general models (II)

Add m new node at each time $t = 0, 1, \dots$ with $m \geq 1$
